ELSEVIER
Discrete Applied Mathematics 113 (2001) 299–302
DISCRETE APPLIED MATHEMATICS

Note

Optimal quantitative group testing on cycles and paths ☆

F.K. Hwang[a] J.S. Lee[b, *]

[a] Department of Applied Mathematics, National Chiao-tung University, Hsin-Chu 300, Taiwan, ROC
[b] Department of Mathematics, National Kaohsiung Normal University, Kaohsiung 802, Taiwan, ROC

Abstract

We determine the minimum number of group tests required to search for a special edge when the graph consists of cycles and paths, generalizing previous results of Aigner on paths and on a simple cycle.

Keywords: Group testing; Search; Graph

1. Introduction

Suppose that we have a set of items containing exactly two defective ones. The problem is to identify them through quantitative group testing [2]. Any subset S of items can be tested, and the feedback $f(s)$ reveals the number of defectives in S, i.e. $f(S)=0,1$ or 2. There are constraints on which pairs of items can be the defective pair, and the constraints can be represented by a graph where the vertex-set is the set of items, and the edge-set is the set of allowed pairs. Thus, the problem can also be viewed as searching for a special edge on a graph $G(V,E)$.

Suppose $|E|=n$. Since each test has three possible feedbacks, $\lceil \log_3 n \rceil$ is the information lower bound on the number of tests required. Aigner [1] proved

Theorem 1. *If G consists of paths, then $\lceil \log_3 n \rceil$ tests suffice.*

Theorem 2. *If G is a cycle and $n<3^t$, then t tests suffice. If $n=3^t$*, then $t+1$ *tests suffice.*

☆ This research is supported by NSC grant 88-2115-M-009-016.
* Corresponding author.

PII: S0166-218X(01)00188-3

In this paper we consider the case that G consists of any number of cycles and paths. We give the minimum number of tests required for such G.

2. Optimal testing

We first prove an upper bound.

Theorem 3. *Suppose G consists of cycles and paths. Then $1 + \lceil \log_3 n \rceil$ tests suffice.*

Proof. If G contain no cycles, then Theorem 3 follows from Theorem 1. If G has m cycles $C_1, C_2, \ldots, C_m$, test $S = \{v_1, v_2, \ldots, v_m\}$, where v_i is an arbitrary vertex on C_i. Suppose $f(S) = 0$. Then the two edges incident to v_i on C_i cannot be special for each $i = 1, 2, \ldots, m$. Therefore C_i is reduced to a path. By Theorem 1, $\lceil \log_3 n \rceil$ more tests suffice. Suppose $f(S) = 1$, then the special edge must be an edge incident to one of the v_i. Again, each C_i is reduced to a path of two edges and Theorem 1 applies. The proof is completed by noting that $f(S)$ cannot be 2 since no edge can be incident to two vertices in S. □

Consider a test S on a graph G. An edge (u, v) will be called an S_i-edge, $i = 0, 1, 2$ if $|\{u, v\} \cap S| = 0, 1, 2$, respectively. Let G_0, G_1, G_2 be the partition of G according to the three feedbacks of S. Then $G_i = \{S_i\text{-}edge\}$ for $i = 0, 1, 2$. A cycle (path) will be called a *mixed cycle (path)* if it contain an S_1-edge. Otherwise it is called a *pure cycle (path)*, or an S_0 (S_2)-cycle if we want to be more specific. We also refer to an edge as *pure* if it is either S_0 or S_2.

Lemma 4. *Let i and j satisfy the conditions $i \geqslant 0$, $j \geqslant 0$ and $i + 2j \leqslant k$, except when $j = 0$, then i is 0 or k. Then there exists a test S on a k-cycle C such that $|S_0| = i$, $|S_1| = 2j$ and $|S_2| = k - i - 2j$.*

Proof. If $j=0$, then either $S \cap C = C$ or $S \cap C = \emptyset$. Otherwise, assign arbitrary $k-i-2j+1$ consecutive vertices to S, and assign the next $i+1$ consecutive vertices to $\bar{S}$ (not in S). The remaining vertices are assigned S or $\bar{S}$ such that S and $\bar{S}$ alternate. □

Lemma 5. *Consider a set P of paths with k total edges. Let i and j satisfy the conditions $i \geqslant 0$, $j \geq 1$ and $i + j \leqslant k$. Then there exists a test S on P such that $|S_0| = i$, $|S_1| = j$ and $|S_2| = k - i - j$.*

Proof. We order the paths such that the k edges (hence all vertices) are linearly ordered. Assign the first $k - i - j$ edges to S_2, meaning their vertices are all in S. Assign the next j edges to S_1, if j is odd or $i = 0$. If j is even and $i > 0$, assign the next $j - 1$ edges to S_1. Furthermore, if there is a change of path during this process, then the vertex starting the new path is in the same set, S or $\bar{S}$, as its preceding vertex. These rules assure that this process ends in an $\bar{S}$-vertex which will start the final assignment of i edges in S_0, meaning all their vertices are in $\bar{S}$. For j even and $i>0$, there is one edge left which will be assigned to S_1, meaning the last vertex is in $\bar{S}$. □

Corollary 6. *A partition $(i,0,k-i)$ is possible if and only if there exists a subset of paths with a total of i edges.*

Let $M(G)$ denote the minimum number of tests required for G.

Theorem 7. *Let G consist only of cycles and paths with n edges in total, where $3^{t-1}<n\leqslant 3^t$. Then $M(G)=t$ except*

(i) *G consists of cycles only and $n=3^t$,*
(ii) *$t=2$ and G contains two cycles,*
(iii) *$t=3$ and G contains seven cycles,*
(iv) *$t=4$ and G contains 26 cycles,*
and $M(G)=t+1$ in the four exception cases.

Proof. *Sufficiency:* The $t\leqslant 2$ case is easily verified. We prove the general $t\geqslant 3$ by induction. It suffices to prove that if G is not one of the exception cases, then there exists a test S where the three feedbacks partition G into G_0, G_1,G_2 with n_0, n_1, n_2 edges, where $n\leq 3^{t-1}$ and G_i is not an exception case for $i=0,1,2$.

Suppose G contains c cycles where $c\leq 3^{t-1}-1$. We consider two cases:

(1) $c<(3^{t-1}-1)/2$. Assign S_1-edges such that the c cycles are all mixed. Suppose the c cycles contain n' edges. By Lemma 4 we can obtain at least $2\lceil(n'-c)/2\rceil$ S_1-edges. Assign $\min\{2\lceil(n'-c)/2\rceil,3^{t-1}-1\}=3^{t-1}-j$ edges to S_1, where $j\geq 1$ is odd . Again by Lemma 4, the pure edges in the c cycles can be divided evenly into S_0 and S_2. Since $3^{t-1}-j\geqslant\lfloor n'/3\rfloor$, so $3^{t-1}-j<\lfloor n/3\rfloor$ implies the existence of paths with a total of more than j edges. By Lemma 5, we can obtain j S_1-edges and divide the other edges evenly into S_0 and S_2. Note that in the case $3^{t-1}-j\geqslant\lfloor n/3\rfloor$, even though no S_1-edge is needed on the paths, some S_1-edges may be forced in the process of dividing the path edges evenly into S_0 and S_2. By Lemma 5, at most 1 S_1-edge needs to be forced. This is alright since $3^{t-1}-j+1\leq 3^{t-1}$.

(2) $c\geqslant(3^{t-1}-1)/2$. We will assign the $(3^{t-1}-1)/2$ largest cycles to be mixed each with two S_1-edges. Let p denote the largest size of the pure cycles. Then $p\leqslant 5$ for otherwise the mixed cycles would have consumed $3(3^{t-1}-1)=3^t-3$ edges and there are not enough edges left for a pure p-cycle with $p\geqslant 6$. Let (e_0,e_2) be a division of edges into the S_0 and S_2 type through assigning the pure cycles into G_0 or G_2. Then there is a division with $|e_0-e_2|\leqslant 5$. For $t\geqslant 3$, there are at least four mixed cycles with 12 pure edges on them. By Lemmas 4 and 5, we can divide these pure edges as well as the pure edges on paths (if any) arbitrarily, i.e. the $n-3^{t-1}$ ($n-(3^{t-1}-1)$ if no paths exist) pure edges can be divided evenly into G_0 and G_2. Therefore $n_i\leqslant 3^{t-1}$ for $i=0,1,2$. Furthermore, the number of cycles in G_0 or G_2 is at most

$$\left\lceil \frac{3^{t-1} - 1 - (3^{t-1} - 1)/2}{2} \right\rceil < 3^{t-2} - 1 \quad \text{for } t \geqslant 5,$$
$$\left\lceil \frac{25 - (3^3 - 1)/2}{2} \right\rceil = 6 \quad \text{for } t = 4,$$
$$\left\lceil \frac{6 - (3^2 - 1)/2}{2} \right\rceil = 1 \quad \text{for } t = 3.$$

Hence they are not exception cases.

That $t + 1$ tests suffice for the exception cases follow from Theorem 3.

Necessity: That t tests are necessary for the nonexception case follows from the information lower bound. We now prove that the exception cases cannot be done in t tests.

(i) Since the number of S_1-edges on a cycle must be even, there is no way to partition 3^t edges on cycles into 3^{t-1}, 3^{t-1} and 3^{t-1}.

(ii) Suppose G contains two cycles. Then the number of S_1-edges on these two cycles must be 2 (it must be even). That means one of the two cycles, of size k, is pure. If $k > 3$, then one more test cannot do it by information lower bound. If $k = 3$, then again one more test cannot do it since it is the exception case (i).

(iii) Suppose G contains seven cycles. Since at most $(3^{3-1} - 1)/2 = 4$ cycles can be mixed, there are at least three pure cycles. Without loss of generality, assume there are two S_0-cycles. Then G_0 contains two cycles and is the exception case (ii), hence it cannot be done in two more tests.

(iv) Suppose G contains 26 cycles. Since at most $(3^{4-1} - 1)/2 = 13$ cycles can be mixed, there are at least thirteen pure cycles. Without loss of generality, assume there are seven S_0-cycles. Then G_0 contains seven cycles and is the exception case (iii), hence it cannot be done in three more tests. □

References

[1] M. Aigner, Search problem on graphs, Discrete Appl. Math. 14 (1986) 215–230.

[2] D.Z. Du, F.K. Hwang, Combinatorial Group Testing and Its Applications, 2nd edition, World Scientific, Singapore, 2000.